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This is a peer-reviewed, post-print (final draft post-refereeing) version of the following published document, The final publication is available at Springer via http://dx.doi.org/10.1007/978-3-319-57186-7_31. and is licensed under All Rights Reserved license:

**Trovati, Marcello, Win, Thu Yein ORCID logoORCID:
<https://orcid.org/0000-0002-4977-0511>, Sun, Quanbin and
Kontonatsios, Georgios (2017) Assessment of Security
Threats via Network Topology Analysis: An Initial
Investigation. Green, Pervasive, and Cloud Computing, 10232.
pp. 416-425. doi:10.1007/978-3-319-57186-7_31**

Official URL: https://doi.org/10.1007/978-3-319-57186-7_31

DOI: http://dx.doi.org/10.1007/978-3-319-57186-7_31

EPrint URI: <https://eprints.glos.ac.uk/id/eprint/5492>

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Assessment of Security Threats via Network Topology Analysis: An Initial Investigation

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Abstract. Computer networks have increasingly been the focus of cyber attack, such as botnets, which have a variety of serious cybersecurity implications. As a consequence, understanding their behaviour is an important step towards the mitigation of such threat. In this paper, we propose a novel method based on network topology to assess the spreading and potential security impact of botnets. Our main motivation is to provide a toolbox to classify and analyse the security threats posed by botnets based on their dynamical and statistical behaviour. This would potentially lead to a better understanding and prediction of cybersecurity issues related to computer networks. Our initial validation shows the potential of our method providing relevant and accurate results.

Keywords: Botnets, Cybersecurity, Network Theory

1 Introduction

Security threats have been steadily increasing due to the emergence of new technology and methodologies, which has led to an expanding research effort to detect and minimise such threats [1, 2]. More specifically, *botnets* due to their unique structure based on distributed communication command patterns across networks, are widely regarded as a serious security issue. In fact, they can successfully carry out surveillance attacks, perform DDoS extortion, general spam, as well as phishing. Furthermore, some of them utilise structured overlay networks, whose lack of centralisation enhance the ability of a botnet to evade detection whilst retaining a good level of robustness with respect to a churn process, where single machines are frequently cleansed [6]. It is estimated that their use has led to malicious activity resulting in a loss of millions of dollars per year [5].

In this paper, we introduce a novel method to assess security threats, based on the dynamical properties associated with networks generated by computer communication. In fact, their topology can provide an insight into specific features exhibited by botnets across computer networks. How connections change, their types and length of communication can provide a deeper and more efficient

Table 1. A selection of network connection flows

Time	Source	Destination	Protocol	Length
161.077519	PcsCompu.b5:b7:19	Broadcast	ARP	60
162.079007	PcsCompu.b5:b7:19	Broadcast	ARP	60
162.079013	PcsCompu.b5:b7:19	Broadcast	ARP	60
162.765245	147.32.84.165	147.32.84.255	NBNS	110
162.765253	147.32.84.165	147.32.84.255	NBNS	110
166.206344	147.32.80.9	147.32.84.165	DNS	503
166.207297	147.32.84.165	74.125.232.195	TCP	62
166.207308	147.32.84.165	74.125.232.195	TCP	62
166.215343	74.125.232.195	147.32.84.165	TCP	62
166.21559	147.32.84.165	74.125.232.195	TCP	60

approach to security threat detection and prediction.

To achieve this, we consider five main parameters: *time*, *source*, *destination*, *protocol*, and *length*. Table 1 depicts a small example of these parameters of the connection flows.

The main motivation is to provide a set of tools to assess the behaviour of host-to-host communication to allow an agile, real-time assessment. In contrast to the the current state of the art approaches, which tend to focus on the different parameters based on whether or not they are present in the collected data, we are aiming to exploit the topology of the network and the probabilistic information related to botnets behaviour. In fact, a dynamical investigation of such networks, can lead to the assessment of the likelihood of the maliciousness of computer communications.

The paper is structured as follows. In Sections 2 and 3 we provide a description of existing technology and theories, and in Section 4 we detail our approach. In Section 5 we discuss the validation process and finally, Section 6 concludes our work and prompts to future research directions.

2 Related Work

In [3], the authors propose a detection method for botnets from large datasets of Netflow data, based on a variety of cloud computing paradigms especially MapReduce for detecting densely interconnected hosts which are potential bot-net members.

BotGrep [5] is a tool to identify peer-to-peer communication structures based on the information about communicating pairs of nodes. This type of P2P detection is defined as a (communication) network, which exploits the spatial relationships in communication traffic. Furthermore, the authors argue that subnetworks with different topological patterns can be partitioned by using random walks, whilst

comparing the relative mixing rates of the P2P subnetwork structure and the rest of the communication network. However, such approach is computationally expensive due to the typical size of such networks.

In [4], an approach based on a Markov chain model is introduced. In particular, botnet infection is modelled to identify behaviour that is likely to be associated with attacks, with a prediction rate over 98%. Another example of the utilisation of Markov chain for intrusion detection system is described in [8], where it is trained on a sequence of audit events. However, these types of approaches allow attack identification but they have limited intrusion prediction. In [9], the authors assess the set of bot lifecycle stages using Markov chains to identify the occurrence of infection. Similar to the previous approach, there is a focus on the identification of infection rather than on any predictive capability.

3 Network Theory

Networks have been extensively used to successfully model many complex systems, and their applications span across a variety of multidisciplinary research fields, ranging from mathematics and computer science, to biology, and the social sciences [12], [13].

Networks are defined by a *node set* $V = \{v_i\}_{i=1}^n$, and the *edge set* $e_{v_i, v_j} \in E$, so that if v_a and $v_b \in V$ are connected, then $e_{v_a, v_b} \in E$ [2]. Note that in this paper, we do not allow self-loops, or in other words, $e_{v_i, v_i} \notin E$.

Scale-free networks, in particular, appear in a numerous contexts, such as the World Wide Web links, biological and social networks [2]. The main property of scale-free networks is based on their node degree distribution, which follows a power law. More specifically, for large values of k , the fraction p_k of nodes in the network having degree k , is modelled as

$$p_k \approx k^{-\gamma} \quad (1)$$

where γ has been empirically shown to be typically in the range $2 < \gamma < 3$ [2]. From Equation 1, it follows that a relatively small number of hubs occur, which define the topological properties of the corresponding networks, as well as the way information spreads across them [15].

An important property of such networks is related to the creation of new nodes over time, which are likely to be connected to existing nodes that are already well connected. Since the connectivity of nodes follows a distribution which is not purely random, the dynamical properties of such networks and their general topological properties can lead to predictive capabilities [13].

4 Description of the Method

In this section we introduce the model whose objective is to understand, assess and predict the type and severity of security threats. As discussed above, the

dynamical properties of networks can provide a useful insight into the system they model. In this paper, we will focus on the following properties:

- The topology of the network, or in other words, the level of connectedness between nodes measured by joining paths, and
- Their dynamical properties.

Loosely speaking, we are interested in the properties exhibited by the single threats and how they change over a specific amount of time.

As defined in Section 3, let $G = G(V, E)$ be a directed network where V is the node set and E is the arc set. The former contains the nodes, and the latter contains the arcs, or directed edges corresponding to requests from the source node to the target node. Let $\deg_{in}^t(v_i)$ and $\deg_{out}^t(v_i)$ be the in and out degrees of the node v_i at a given time t , that is the number of connection into and out of it, respectively. We then define the *maliciousness* of a node v_i at the time t , as

$$P_M^t(v_i)_{in} = \frac{|\deg_{in}^{M,t}(v_i)|}{|\deg_{in}^t(v_i)|} \quad (2)$$

or

$$P_M^t(v_i)_{out} = \frac{|\deg_{out}^{M,t}(v_i)|}{|\deg_{in}^t(v_i)|}, \quad (3)$$

where $|\deg_{in}^{M,t}(v_i)|$ and $|\deg_{out}^{M,t}(v_i)|$ are the number of malicious connections into or out of v_i , at a given time t .

For a time t and an arc $e_{v_i, v_j} \in E$, define its *weight* as

$$w_t(v_i, v_j) = f_t(r, p), \quad (4)$$

where $f_t(r, p)$ is a function of the *length of time* of a request r and the *number of request protocols* p from v_i and v_j . In this paper, we define

$$f_t(r, p) = \frac{1}{2}(w_r^t + w_p^t), \quad (5)$$

where w_r^t and w_p^t are the length of the time and the number of protocols of different requests, respectively.

We then define the probability of a malicious request at the time t from v_j to v_i as

$$P_M^t(v_i, v_j) = \frac{1}{3}(P_M^t(v_i)_{in} + P_M^t(v_j)_{out} + w_t(v_i, v_j)). \quad (6)$$

In order to consider the dynamics of this model, we assume that new requests arise according to time snapshots $t = 1, \dots, T$. Let

$$\delta_T(v_i, v_j) = P_M^t(v_i, v_j) - P_M^{t-1}(v_i, v_j) \quad (7)$$

and define

$$\Delta_T(v_i, v_j) = \frac{1}{T-1} \sum_{t=2}^T \delta_T(v_i, v_j). \quad (8)$$

Finally, let the *probability of maliciousness* as

$$\tilde{P}_M^T(v_i, v_j) = \min \{ \max \{ \Delta_T(v_i, v_j), 0 \}, 1 \}. \quad (9)$$

Note the above equation can be extended to assess the (average) probability of malicious attacks from a set of nodes \tilde{V} on a specific node as

$$\tilde{P}_M^T(v_i) = \frac{1}{|\tilde{V}|} \sum_{\tilde{v} \in \tilde{V}} \tilde{P}_M(v_i, \tilde{v}) \quad (10)$$

for $e_{v_i, \tilde{v}} \in E$ and \tilde{v} is a node in \tilde{V} .

Algorithms 1 and 2 show the implementation of the above approach.

Algorithm 1 Evaluation of $\tilde{P}_M^T(v_i, v_j)$

- 1: Let $t = 0$
 - 2: Determine $P_M^{t=0}(v_i, v_j)$
 - 3: **for** $t = 1, \dots, T$ **do**
 - 4: Find $\Delta_T(v_i, v_j)$ and $\tilde{P}_M^T(v_i, v_j)$
 - 5: **end for**
 - 6: **return** $\tilde{P}_M^T(v_i, v_j)$
-

Algorithm 2 Evaluation of malicious attacks on node v_i

- 1: Let $t = 0$
 - 2: **for** $\tilde{v} \in V \setminus v_i$ **do**
 - 3: Determine $P_M^{t=0}(v_i, \tilde{v})$
 - 4: **for** $t = 1, \dots, T$ **do**
 - 5: Find $\Delta_t(v_i, \tilde{v})$ and $\tilde{P}_M^T(v_i, \tilde{v})$
 - 6: **end for**
 - 7: **end for**
 - 8: **return** $\tilde{P}_M^T(v_i, \tilde{v})$
-

As discussed in Section 3, if the network G follows a scale-free structure, new arcs are likely to be added to highly connected nodes. As a consequence, Equations 2 and 3 can be modified to incorporate this property. Recall that the fraction of nodes p_k with degree k is

$$p_k \approx k^{-\gamma}.$$

For a scale-free network G , we then assume that

$$P_M^t(v_i, v_j) = \frac{1}{2}(\deg(v_i)^{-\gamma} + w_t(v_i, v_j)). \quad (11)$$

Note that in this case, we are considering the overall degree of the destination node v_i , rather than distinguishing between the in and out degree values of the source and destination nodes. Although we are providing fewer parameters in the model above, compared to Equation 6, the initial validation appears to support the claim that (11) indeed provides good modelling capabilities.

The dynamics described by Equation 9 can be used to provide some level of prediction of the number of malicious attacks. In this paper, we assume that the trend of $\tilde{P}_M^T(v_i, v_j)$ can give an insight into a “near future” behaviour of the communications from v_j to v_i . In particular, we shall assume that $\tilde{P}_M^T(v_i, v_j) \approx \tilde{P}_M^{T+1}(v_i, v_j)$, or in other words, they exhibit a similar trend. We acknowledge this is a simplistic approach as it does not consider potential variations that could occur. However, our initial validation seems again to support the above. In future research, we are aiming to fully investigate and extend the predictive properties of our approach by fully analysing the topology of a large set of communication networks.

5 Results

In this section, we will discuss the validation process, which was based on the publicly available datasets offered by the Malware Capture Facility Project [11]. More specifically, we used the CTU-MALWARE-CAPTURE-BOTNET-42 dataset, which contains relevant data generated by a Neris botnet. It used an HTTP based C&C channel, and all the actions performed by the botnet were communicated via C&C channels containing specific “click-fraud” spam based on advertisement services.

This was subsequently preprocessed via WireShark [10] to capture all the parameters relevant to our approach.

A directed network $G = G(V, E)$ was defined, where the node-set V contains the source and destination IPs mutually linked by a request. In particular, we had

- Number of nodes: 4247
- Number of arcs: 6588
- Average in and out degree: 1.5512

Figures 1 and 2 show the degree distribution of the network G , which indicates the existence of few highly connected hubs. Note that this behaviour is similar to scale-free networks, as described by Equation 1. In [14], a method to topologically reduce complex networks is discussed. When such method is applied to the network G , a value of $\gamma = 1.9$ is determined. As discussed in Section 3, for many complex systems γ is usually within the range $2 < \gamma < 3$, suggesting that the dataset used for the validation exhibits properties similar to

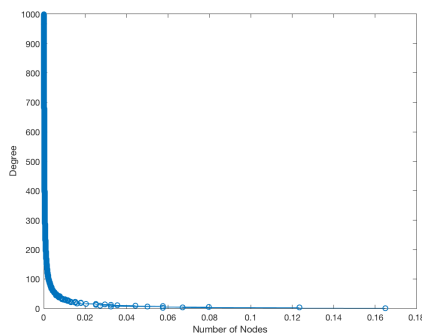


Fig. 1. The degree distribution of the network G .

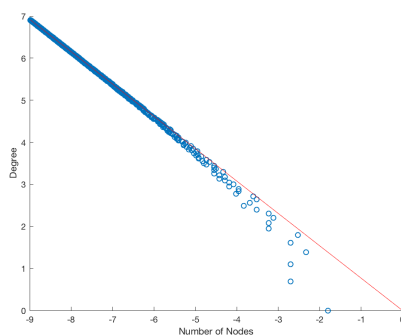


Fig. 2. The log node degree distribution of the network G depicted in Figure 1, which is compared with a (theoretical) scale-free network with $\gamma = 1.9$. As it can be clearly seen, this is a good approximation of the node degree distribution of G .

many other systems from across various contexts. As discussed above, we are aiming to widen our investigation to a large set of malware botnet datasets to fully assess whether such behaviour can be indeed generalised.

In order to evaluate our approach, we trained the parameters of Equation 5 on approximately 2000 malicious requests. First of all, we noticed that over 95% of the malicious requests had a TCP protocol, and among them we detected two main clusters for time length values in the interval $[0, 70]$ and $[950, 1400]$, as depicted in Figure 3.

Therefore, we assumed that for malicious requests, $w_t = w_p = 0.6$ if the time length is within those intervals and the protocol is TCP, and $w_t = w_p = 0.2$ otherwise.

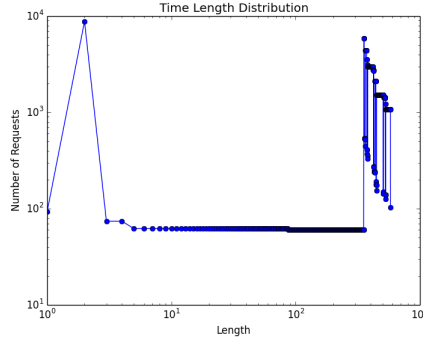


Fig. 3. The distribution of the time length requests.

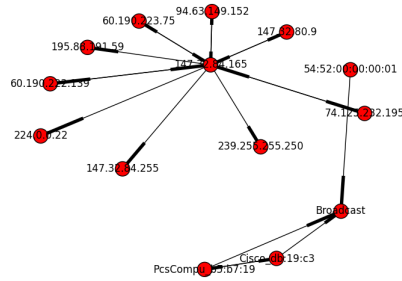


Fig. 4. A sub-network generated by the dataset described in Section 5.

We subsequently considered the dynamics of the system, by analysing batches of approximately 300 requests per time iterations, and we assumed that a malicious request from v_j to v_i is associated with $\tilde{P}_M^T(v_i, v_j) > 0.7$. The analysis of the data produced that 71% of the malicious requests had indeed a $\tilde{P}_M(v_i) > 0.7$. Figures 4, 5 and 6 depict a small proportion of the network created in the first three iterations on the process. Furthermore, Figure 6 also shows the malicious requests, which are depicted in red.

We subsequently evaluated the model defined by Equation 11. In this case, $\tilde{P}_M(v_i, v_j) > 0.7$ for approximately 61% of the malicious requests. This decrease in accuracy was indeed expected due to the more general scope of the model, as discussed above.

Finally, we evaluated the level of prediction associated with our model, and we considered approximately 200 pairs of nodes exchanging request. Approximately 59% of the malicious requests exhibited the same trend $\tilde{P}_M^T(v_i, v_j) \approx \tilde{P}_M^{T+1}(v_i, v_j)$, and we noted this was particularly the case for larger values of T , as expected.

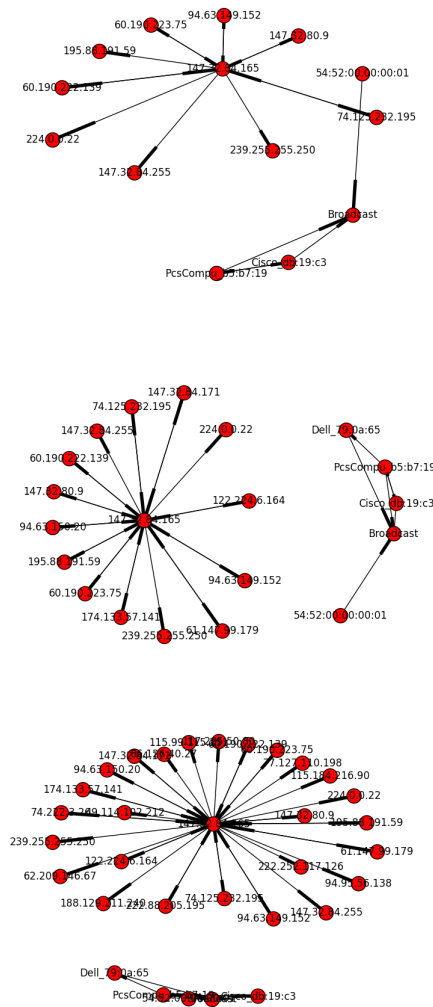


Fig. 5. A selection of the sub-networks generated by some time iterations on the dataset described in Section 5.

6 Conclusion

In this paper, we have discussed a method to assess and predict the malicious connection requests in terms of bonets. As indicated by the validation shows, this approach shows potential in providing a robust method to detect and predict malicious request activity. However, this is still at its infancy and in future research we are aiming to extend our investigation to consider more parameters

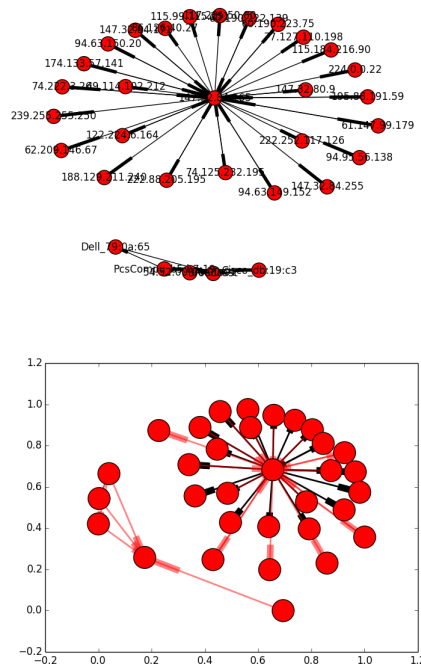


Fig. 6. The sub-networks generated by the fourth iteration, where the second figure highlights the malicious connections.

and create a more comprehensive model. In particular, a full investigation of networks generated by such requests will require a deeper understanding of the topological properties of such networks to ensure a more comprehensive and accurate analysis, which will provide a robust, accurate and computationally effective approach.

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